

A Simple Algorithm to Calculate $S(n)$

by John C. McCarthy

Introduction

This short paper first outlines an "obvious" algorithm for calculating $S(n)$ (the smallest integer m such that $m!$ is divisible by n). Doubtless, there exist more subtle and efficient algorithms. I hope some readers will devise these and enlighten me concerning them through this journal.

This is followed by a small scale investigation of the efficiency of the algorithm.

Then there is a short discussion of a simple way of reducing the space required for storage of all $S(n)$ for ranges of n . The storage space required for $S(n)$ for all n which my routines can handle is considered.

Heavily commented listings of an implementation of the algorithm in "C", sample output and timing data are included to help illustrate the algorithm.

The Algorithm

The algorithm is described in detail at the start of the header file "S(n).H". Together with "S(n).C", this forms all the code necessary to implement the algorithm. Note that, for the $S(n)$ function to work correctly, the function `make_primes()` must first be called from the main program.

The code for printing $S(n)$ and timing the routines has been omitted. These activities are both implementation specific and easily done. They are therefore left as an exercise for the interested reader.

The algorithm hinges on finding the prime factors of n . Improvements on how this is done will most benefit its efficiency.

To be practical, the given implementation of the algorithm only works for $0 < n < 2^{32}$. However, the algorithm is generally applicable to any non-null integer.

Tables of $S(n)$, constructed using the routines of "S(n).C", for the largest 2000 permitted n are included.¹ My paging routines are rather elaborate. Using them (without printing!), it took 2.4 hours to discover that 3,745,708 pages, as tightly packed as those shown, would be required to print $S(n)$ for all $0 < n < 2^{32}$.

Efficiency of the Algorithm

In a letter to R. Muller (about computing the Smarandache Function, July 19, 1993), Ian Parberry (editor of <SIGACT News>,

¹ For the smallest 4800 numbers, see Istedt's table (pp. 43-50) of this current journal.

Denton, Texas) expressed that one can immediately find an algorithm that computes $S(n)$ in $O(n \log n / \log \log n)$ time ('A Brief History of the "Smarandache Function"' by Dr. Constantin Dumitrescu, Department of Mathematics, University of Craiova, Romania). Disappointingly, a little analysis of the accompanying timing data on my TI85 advanced scientific calculator reveals that my algorithm is somewhat worse than this.

Trying to fit the version 2 timing data to various $O(f(n))$, I obtained the following results ($x=3355443200$ and $10(O(x+99)-O(x-100))$ is calculated for comparison with the last entry of the version 1 timing data):

$O(f(n))$	Correlation Coefficient	$O(2^{3^2}-1)$ (years)	$10(O(x+99)-O(x-100))$ (milliseconds)
$O(n)$	0.9928879	0.6092	8909
$O(n \log n / \log \log n)$	0.9944006	0.7906	11827
$O(n\sqrt{n})$	0.9997756	24.2	469178

$O(n\sqrt{n})$ fits the version 2 timing data best, although the time it predicts for the last entry of the version 1 timing data is almost 3 times too large. Hence, I assume the time complexity of my algorithm is a little better than $O(n\sqrt{n})$.

As a rough upper limit on the time my program (on my 20MHz 368DX PC) would take to calculate $S(n)$ for all $0 < n < 2^{3^2}$, let us assume that every value of n requires as much time as each n in the range of the last entry of the version 1 timing data ($= 159111/199/10 = 79.9553$ ms). In this "worst case", it would take 10.882 years. $O(n\sqrt{n})$ time complexity predicts more than twice this value, which is a measure of how pessimistic it is.

I would welcome a more rigorous analysis of the time complexity of my algorithm as I presently lack the necessary expertise.

Simple Compression of Stored $S(n)$

Without compression, each $S(n)$ would be stored as a 32-bit (= 4 bytes) value. Hence 2^{3^2} bytes (= 16 Gigabytes) would be required to store $S(n)$ for all $0 < n < 2^{3^2}$.

This requirement can be reduced considerably if we use the high bit of each byte of each value to indicate if it is the last byte of the value. If the bit is set it means that further byte(s) are required and if it is reset it means that the byte is the last byte of the current value. This means that only 7 bits of each byte are used to form the numerical part of the value. Assuming that, as with Intel format, the values are stored low-'byte' (actually 7 bits) first, here are some examples:

- i) 127 requires seven bits and so just one byte (with high bit reset to indicate no further bytes).
- ii) 16,000 requires 14 bits. So it is stored as two bytes. The

first is 0 (16,000 mod 128) + 128 (to set the high bit indicating there is more to come). The second is 125 (16,000 div 128) (with high bit reset to indicate no further bytes). This reads simply as 0 (with more to follow) + 128*125 (no more to follow).

- iii) A number stored as the three bytes 57+128, 93+128 and 125+0 would similarly represent:
 $57 + 93*128 + 125*128*128 = 2,059,961.$

The largest numbers that can be represented by a given number of bytes is thus as follows:

- 1 byte can code up to $2^7-1 = 127.$
- 2 bytes can code up to $2^{14}-1 = 16,383.$
- 3 bytes can code up to $2^{21}-1 = 2,097,151.$
- 4 bytes can code up to $2^{28}-1 = 268,435,455.$
- 5 bytes can code up to $2^{35}-1 = 34,359,738,355$ (or 8 times the largest unsigned long).

For small values of n, the savings are considerable (400%). However, even large n often have small S(n).

Using this technique to compress all S(n) calculated for some ranges of n (each range was also stored), I obtained the following results:

range of n	compression	time taken (seconds)	size	size after pkzip
1 -10,000	without	4.5	40,008	19,836
	with	4.7	15,749	15,267
2,147,478,648 -2,147,488,647	without	827.3	40,008	33,729
	with	842.4	33,541	30,836
4,294,957,296 -4,294,967,295	without	1,066.2	40,008	34,320
	with	1,085.1	34,330	31,634

The results indicate that this compression is a little better than pkzip's (a commercial file compression utility). Application of pkzip to a pre-compressed file also gives a slight improvement.

Assuming that the savings shown for the middle range of 10,000 n are the average of all ranges of 10,000 n, using my compression together with that of pkzip would permit storage of S(n) for all $0 < n < 2^{32}$ in about $3.0836 * 2^{32} = 12.3344$ Gigabytes. So look out for sets of 19 CD-ROMs with all your favourite numbers on them!

21st November 1993

```
/* (c).1993.11.13.John.C.McCarthy
   "S(n).h"
```

Example Implementation of A Simple Algorithm to Calculate S(n),
The Smarandache Function:

Because there are more people familiar with C than with C++, this module has been written entirely in C (apart from "//" style comments). The module was compiled using Borland C++ version 3.1.

For efficiency, n is constrained to the limits of an unsigned long. Hence, $0 \leq n \leq 2^{32} - 1$ ($= 4,294,967,295$). ("^" represents exponentiation). Although catering for n of vast magnitude is possible, it imposes heavy storage and processing overheads. The range of an unsigned long therefore seems a reasonable compromise.

The algorithm depends on the most elementary properties of S(n):

1) Calculate the STANDARD FORM (SF) of n:

In SF: $n = +/- (p_1^{a_1}) * (p_2^{a_2}) * \dots * (p_r^{a_r})$ where p_1, p_2, \dots, p_r denote the distinct prime factors of n and a_1, a_2, \dots, a_r are their respective multiplicities.

2) $S(n) = \max[S(p_1^{a_1}), \dots, S(p_r^{a_r})]$.

3) $S(p^a)$, where p is prime, is given by:

3.1) $a \leq p \implies S(p^a) = p^a$.

3.2) $a > p \implies S(p^a) = x < p^a$. In this case, fortunately rare, x is the smallest integer such that p appears as a factor in the list of all integers > 1 and $\leq x$ at least a times. Let the no. of times p appears as a factor in the list of all integers > 1 and $\leq y$ be $f(y, p)$. Then:

$f(y, p) = \sum [\text{int}(y/(p^i))] \text{ for } i > 0 \text{ while } y \geq (p^i)$.

Hence, x is the smallest integer such that $f(x, p) \geq a$.

Note that between successive integer multiples of p there are no integers which have p as a factor. The trick here is to look for the largest multiple of p (call it c), such that $f(p^*c, p) \leq a$ (so that $x = p^*c$, if $f(p^*c, p) = a$, else $x = p^*(c+1)$):

3.2.1) $c = a - 2$ (largest possibility for c since $f(p^*(a-1), p) \geq a$ when $a > p$ (Note: $f(p^*(a-1), p) = a$ is not sought for slight performance gain)).

3.2.2) $z = f(p^*c, p)$.

3.2.3) While($z > a$):

3.2.3.1) $d = \text{no. of times } p \text{ appears as a factor of } p^*c$
 $= (\text{no. of times } p \text{ appears as a factor of } c) + 1$.

3.2.3.2) $c = c - 1$ (next largest possibility for c).

3.2.3.3) $z = z - d$ ($= f(p^*c, p)$).

3.2.4) If($z < a$), $x = p^*(c+1)$.

3.2.5) Else $x = p^*c$.

To calculate the prime factors of all 32-bit n requires use only of primes $< (2^{16})$ (i.e. all primes expressible as an unsigned short integer). This is because any factor of n remaining after division of n by all its prime factors $< (2^{16})$ is simply a prime. Since there are only 6542 16-bit primes, the program first creates a list of these (which only takes about 4 seconds on my 20 MHz 386DX PC) so that they never have to be recalculated, thus saving much time.

*/

```

#define PRIMES16 6542    // The number of 16-bit primes
#define MAX_SFK 9  /* max. distinct primes in the SF of n. The smallest
    number with more than 9 distinct primes is the product of the 10 smallest
    primes (= 6,469,693,230), which is substantially more than the largest
    integer expressible as an unsigned long. Hence, 9 distinct primes are
    more than ample.
*/

typedef unsigned long u_long;
typedef unsigned int  u_int;
typedef enum {false, true} boolean;

struct SF_struct {
    int      sfk;           // no. of distinct primes
    u_long  sfp[MAX_SFK];  // the distinct primes
    int      sfa[MAX_SFK]; // respective multiplicities
};

extern u_int prime[PRIMES16+1]; // list of all 16-bit primes
                                // plus terminating zero.

void make_primes(void); // construct list of all 16-bit primes (prime[]).
                        // Must be called before calls to getSF() or S().
void getSF(u_long n, struct SF_struct *SF); // calc. SF of n and store in SF
u_long S(u_long n); // calc. S(n)
u_long Spa(u_long p, int a); // calc. S(p^a) where p is prime
int f(int x, int p); /* the number of times the prime p appears as a factor
    in the integers from 1 to x inclusive. This function is only called from
    Spa(p, a) when a>p with x=p*(a-2) (refer to item (3) of algorithm outline
    above). Max value of (a) occurs when p is a minimum, n is a maximum and
    (p^a)=n. So, (2^max(a))=max(n)=(2^32)-1. Hence max(a)<32. So, x<60
    when (a) is at its max. Max value of p (and x) occurs when a=p+1 and
    (p^a)=max(n). So, max(p)^(max(p)+1)=(2^32)-1. The upshot is that
    max(p)=9 when a=10. Hence, max(x)=72. This explains why it is safe for
    x, p and the return value of f(x,p) to be passed as ints.
*/

```

S(n).C

```
/* (c).1993.11.13.John.C.McCarthy
   "S(n).c"
```

Example Implementation of A Simple Algorithm to Calculate S(n),
The Smarandache Function:

This is the code for the module. Refer to "S(n).h" for details.

```
*/
#include "S(n).h"

u_int prime[PRIMES16+1]; // allocate storage for list of all 16-bit primes
                        // plus terminating zero.

void make_primes(void)
{
    u_int *pp; // ptr to last prime so far of prime list
    u_int *tp; // ptr to current test prime
    u_int p;   // number being tested for primality

    pp=prime; // point to start of prime list
    *pp=2;    // set first prime to 2
    *++pp=3;  // set second prime to 3
    p=5;     // next possible prime. N.B. p is kept odd so that trial
            // division by 2 is unnecessary.
    while(true) { // infinite loop!:
        tp=prime+1; // point to first odd test prime
        // whilst test prime <= sqrt(p):
        while(((long) *tp)*(*tp)<=p) {
            if(!(p%*tp)) { // If current test prime divides (is factor of) p:
                p+=2; // try next odd number
                if(p<*pp) { // done when p overflows:
                    *++pp=0; // terminate list
                    return;
                }
                tp=prime+1; // point to first odd test prime
            }
            else ++tp; // Else point to next test prime
        }
        // no prime <= sqrt(p) divides p so p must be prime:
        *++pp=p; // so store it next in the list
        p+=2; // try next odd number
        if(p<*pp) { // done when p overflows:
            *++pp=0; // terminate list
            return;
        }
    }
}
}
```

S(n).C

```

void getSF(u_long n, struct SF_struct *SF)
{
    u_int *pp; // ptr to current prime
    u_long r; // 'residue' of n remaining for factoring

    SF->sfk=0; // no. of distinct prime factors discovered
    r=n;
    pp=prime; // point to start of prime list

    // whilst current prime <= sqrt(r) and prime list not exhausted:
    while(((long) *pp)*(*pp)<=r && *pp) {
        if(!(r%*pp)) { // if current prime is a factor of r:
            SF->sfp[SF->sfk]=*pp; // store current prime as next prime of SF
            SF->sfa[SF->sfk]=1; // set its multiplicity to 1
            r/=*pp; // 'divide out' current prime
            while(!(r%*pp)) { // while current prime factors r:
                SF->sfa[SF->sfk]++; // increment multiplicity
                r/=*pp; // 'divide out' current prime
            }
            SF->sfk++; // increment count of distinct prime factors
        }
        ++pp; // next prime
    }

    if(n>1) { // If n contains prime > 2^16:
        SF->sfp[SF->sfk]=r; // store it as last prime of SF
        SF->sfa[SF->sfk]=1; // set its multiplicity to 1
        SF->sfk++; // increment count of distinct prime factors
    }
}

```

S(n).C

```

u_long S(u_long n)
{
    struct SF_struct SF; // to store SF of n
    int sfi; // index of current term of SF of n
    u_long Sn; // current guess at S(n)
    u_long x; // S(current term of SF of n) where it might exceed
                // current value of Sn.

    if(n==1) return 0; // special case

    getSF(n, &SF); // calc. and store SF of n

    // First guess at S(n) is S(p^a), where p is the largest prime in the SF
    // of n and a is its multiplicity. This pre-empts the calculation of S(p^a)
    // for the remaining terms where, as is likely, p^a for these terms is <=
    // this initial guess (since S(p^a) <= p^a always):
    sfi=SF.sfk-1;
    Sn=Spa(SF.sfp[sfil],SF.sfa[sfil]);

    while(sfi>0) { // while more term(s):
        sfi--; // next term
        if(SF.sfp[sfil]*SF.sfa[sfil]>Sn) { // if this term may have larger S(p^a):
            x=Spa(SF.sfp[sfil],SF.sfa[sfil]); // calc. it
            if(x>Sn) Sn=x; // if new max., update Sn with it
        }
    }
    return Sn; // That's all folks!
}

u_long Spa(u_long p, int a)
{
    // Refer to item 3) of the algorithm description in S(n).h.
    int c; // largest multiple of p such that f(p*c, p) <= a (eventually!)
    int z; // f(p*c, p)
    int m; // used to calc. no. of times p appears as factor of c

    if(a<=p) return p*a;

    c=a-2;
    z=f(p*c, p);
    while(z>a) {
        // d in items 3.2.3.1) and 3.2.3.3) of algorithm description is implicit
        // here:
        z--;
        m=c--;
        while(!(m%p)) { // while p divides m:
            z--;
            m/=p; // 'divide out' factor of p from m
        }
    }
    if(z<a) return p*(c+1);
    else return p*c;
}

```


S(n).C

```
int f(int x, int p)
{
    int k=0; // count of appearance of prime p as a factor in the integers
            // from 1 to x.
    int xdp; // successive divisions of x by p

    xdp=x/p;
    while(xdp>0) {
        k+=xdp;
        xdp/=p;
    }
    return k;
}
```

Table with 10 columns labeled 0 through 9. Each row contains 10 numerical values. The first column values range from 4294965296 to 4294966406. The other columns contain various integers, some with leading zeros, representing the function values S(n).

S of n plus:										
n\	0	1	2	3	4	5	6	7	8	9
4294966416	1877	813566631	48677	8088449	251	30899039	4804213	390451493	7213	57266219
4294966418	1564081	4294966427	7017919	95569	2939	204522211	577	389449	12541	636763
4294966436	37571	17874759	809	15913	261251	4294966441	614093	1431655481	41539	656221
4294966446	3911827	4294966447	1172207	1137137	965181	2669339	78301	3929521	199	5039
4294966456	26881	284827	6587227	172787	202021	19611719	443	99882941	2465513	858933233
4294966466	306783319	130150499	2848121	17970571	143165549	12763	8468323	1949	2147483237	1381
4294966476	38891	4294966477	3308911	7001	681	330382037	321143	2176871	1073741621	2951867
4294966486	89273653	613566641	178956937	3455323	429496649	159072833	56512717	6806603	7866239	858933299
4294966496	7895159	49367431	2147483249	4999961	86787	813566643	5413	1431655501	536870813	6151
4294966506	234487	416623	1153	477218501	81083	1868189	443	390487	1256573	26407
4294966516	306871	49691	1951	124193	8259551	345727	3984199	9565627	10631	7469507
4294966526	26881	284827	1677723	61165647	443237	1210873	443	32611	29417579	40123
4294966536	8521759	482093	2147483269	18088017	398947	2307881	715827757	61356649	1493713	56509
4294966546	165191021	252645091	16193	3023	12271333	1431655517	536870819	4294966553	238609253	32443
4294966556	73883	1129957	8910719	330382043	64373	119083	52377641	617359	530767	16207421
4294966566	34781	65563	463219	511123	201547	25577	5107	186163	648199	364753
4294966576	1405421	293	37693	138547309	214748329	21577	22139003	4294966583	6170929	9439487
4294966586	58040089	75350291	133189	12739	47721851	4294966591	222933	21387993	298303	179593
4294966596	383	4441	9717119	88667	32587	465781	20368	208261	1073741679	28633107
4294966606	15661	58835159	29826157	22253713	1523	110127349	5088523	7993	908111	49053
4294966616	870131	105323	764501	4294966619	6551	1693	21262211	92347	11719	16901
4294966626	79536419	45751	12064513	23469763	1297573	312839	719	8597	306783331	2219621
4294966636	873871	330382049	7639	4294966639	5821	15671	74051149	14549	153151	4567
4294966646	2147483323	1281697	1080223	284831	2202547	4294966651	1073741663	799366	195225757	15877
4294966656	17449	4294966657	1121849	133361	86627	4294966661	20063	756023	63997	26038101
4294966666	556487	429496667	452239	214309	42949667	1237	865907	1553389	715827779	120223
4294966676	757	8251771	4994793	6053	2482200	86447	12030719	7993	908111	49053
4294966686	17959	390451517	2383	36709117	21401	109451	357913891	54366667	1028981	286331113
4294966696	536870837	16582883	21691751	109961	87119	75350293	38569	138547313	12782639	218629
4294966706	2147483353	975907	4492643	16976153	143165557	2280917	536870839	1431655571	701	2111
4294966716	119304631	3607	20164163	431	2194	482689	69371	20259277	9271	4241
4294966726	52377643	64103981	13765919	208889	144583	84349	397	889043	4079	924643
4294966736	1820687	30460757	11864549	32292983	7156879	4123	195225769	55949	3509	858933249
4294966746	115811	383	3853	143165549	17179867	4141723	14913079	68483	79841	65297
4294966756	1073741689	5811863	1299143	198391	4788	89417	46411	15176581	756689	27709463
4294966766	321047	191	27751	4294966769	47721853	116080183	107021	24265349	306783341	15618061
4294966776	1844917	9133	99223	477218531	16519103	2387419	715827797	6936557	178481	154523
4294966786	1490273	148102303	34019	85199	7459	562979	536870849	330382061	10683997	78311
4294966796	1789	43383503	2147483399	17821439	3579139	93229	93503	1431655601	4481	317323
4294966806	761	99882949	32173	204522239	42949681	1072937	32779	4294966813	8021	86129
4294966816	3257	1516049	359171	1768749	214748341	1389958	26581	290789	26679	65297
4294966826	29417581	13229	9502139	4294966829	35081	91382273	349981	477218537	27059	858933367
4294966836	48623	613566691	15449521	88609	349753	38351	20347	19259941	275389	1579
4294966846	2147483423	168737689	14983	131433	2281	6299	983	782183	4137733	65557
4294966856	511793	82559	761249	3731509	23860927	902873	1030957	35383	747751	4517
4294966866	191449	309399	507341717	2905	1298	330382067	1348511	1387	13177	2290649
4294966876	1073741719	4294966877	9767639	917393	26843543	143165527	5804009	138053	271307	221219
4294966886	306783349	338213	23633	1931	3491843	8311	1073741723	911	50647	9651611
4294966896	1489	330382069	1246363	154423	6135687	113783	715827817	59791	22193	95443709
4294966906	5639	178207	357913909	4294966909	2029	1431655637	70051	13687	2797	858933383
4294966916	7307	10928689	82339	641327	1154561	36092159	8555711	453199	3271	36191
4294966926	715827821	4294966927	572357	8263	5436667	104755291	57943	12744709	35204647	21839
4294966936	18121	2403451	42107519	188737693	19522577	22453	306783353	4294966943	9151	858933389
4294966946	20069339	5986191	56512723	47197438	5923	981	536870859	1431655651	2147483477	683
4294966956	1381907	353699	23087	397351	56497	447439	217643	8765238	1073741741	393583
4294966966	70067	226050893	59652319	97189	7621	1431655657	449	390451543	715827829	12113
4294966976	67108659	163487	2147483489	260791	947	4294966981	20297	1431655661	78341	41177
4294966986	12558383	569851	602887	1588963	98081	3427747	323027	18341	2147483497	193
4294966996	119159	4294966997	26107	16330673	4294967	110127359	4106087	554977	13256071	6458597
4294967006	11483469	65519	5835553	1597	143165567	22453	21913097	375467	12707003	858933403
4294967016	36149	795217	14412841	150058	180007	365747	5549053	17401	14128181	10141
4294967026	23917	18541	32537629	4294967029	11608019	5683	536870879	3764213	3392349	858933407
4294967036	2445881	1413283	8355967	20549	241	10710641	96703	9851	34636831	11767033
4294967046	18111	39163	141319	477218561	5009	21582749	357913921	37201	2879	1228889
4294967056	5347	252645121	83459	1755197	214748353	7121	4007	768193	564533	19976591
4294967066	15761	132967	2599859	2408843	143165589	65551	12201811	16088019	1779191	791699
4294967076	119304641	23761	5250571	110127361	107374177	417839	102281121	390451553	342283	12917
4294967086	2147483543	4294967087	1132639	613566727	429496709	21467	9209	11332367	21691753	393303
4294967096	7669581	143165569	2147483549	96553	362	181967	66977	68174081	5157	9199
4294967106	1583	2135737	2609	1431655703	1252177	4294967111	182423	88463	675097	266331141
4294967116	63841	10058471	754297	15053	236507	102871	12859183	17321	81031	14753
4294967126	2147483563	21841	85639	161471	463319	12899	1073741783	80021	186649	129347
4294967136	13147	138547327	1467863	159072857	214748357	1163	2476913	4294967143	7321	576119
4294967146	93368851	13990121	2251031	2044249	85899343	77249	16381	33816639	70653	515293
4294967156	1073741789	1932059	2147483579	5156023	647	4294967161	34919	4886197	381707	858933433
4294967166	1262483	3229	1801	179383	184571	651839	357913931	9739	2147483587	7283
4294967176	87481	104755297	317159	62929	34981	3583	195225781	14767	6197	44171
4294967186	152293	204522247	5449	4294967189	143165573	82791	23342213	29333	7134497	45210181
4294967196	2521	4294967197	1296007	46182443	31033	11576731	238609289	26029	97612891	286331147
4294967206	1809169	91382281	18539	488008	25264513	477218579	292493	75167	85981	13109
4294967216	3780781	1431655739	52377649	11959	25411	39403369	813749	3089	1726273	231223
4294967226	22097	252645131	1							

Timing of S(n).c Module for Calculation of Smarandache Function, version 1

Time taken to calculate S(n) depends on how easy it is to factor n. Less time is required if n has "small" prime factors. So, in the following table, the values of n shown are the mid-points of ranges (n-99 thru n+99). Times shown are for calculating S(n) for all integers in each range 10 times over:

n	time (ms)
100	268
200	308
400	345
800	387
1600	432
3200	490
6400	571
12800	661
25600	766
51200	919
102400	2450
204800	4036
409600	5670
819200	7977
1638400	10423
3276800	13004
6553600	16302
13107200	23438
26214400	29642
52428800	37011
104857600	50330
209715200	62363
419430400	77888
838860800	108179
1677721600	158480
3355443200	159111

Timing of S(n).c Module for Calculation of Smarandache Function, version 2

"Time to n" is the time taken to calculate S(n) for all n <= that shown.
 "Time add." is the time taken to calculate S(n) for all n > previous n and
 <= current n. All times are in milliseconds (as per version 1):

n	Time to n	Time add.
50000	18223	18223
100000	66763	48540
150000	139191	72428
200000	229634	90443
250000	335252	105618
300000	452539	117287
350000	579419	126880
400000	715146	135727
450000	859963	144816
500000	1012335	152372
550000	1171221	158886
600000	1336899	165678
650000	1508825	171927
700000	1686808	177983
750000	1870023	183215
800000	2058983	188961
850000	2252457	193473
900000	2450892	198435
950000	2653620	202728
1000000	2860734	207115
1050000	3072049	211314
1100000	3288502	216454
1150000	3509106	220603
1200000	3733965	224860
1250000	3962171	228206
1300000	4194158	231987
1350000	4429331	235173
1400000	4668560	239229
1450000	4910513	241953
1500000	5155601	245088
1550000	5404652	249051
1600000	5656512	251859
1650000	5911306	254794
1700000	6169686	258380
1750000	6431383	261697
1800000	6696172	264789
1850000	6963206	267034
1900000	7232974	269768
1950000	7505412	272438
2000000	7779763	274351
2050000	8056579	276816
2100000	8336442	279863
2150000	8620053	283611
2200000	8905641	285588
2250000	9194727	289086
2300000	9486449	291722
2350000	9780105	293655
2400000	10076920	296815
2450000	10375202	298282
2500000	10676383	301181

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