# FLORENTIN SMARANDACHE <br> An Integer Number <br> Algorithm To Solve Linear <br> Equations 

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## AN INTEGER NUMBER ALGORITHM TO SOLVE LINEAR EQUATIONS

An algorithm is given that ascertains whether a linear equation has integer number solutions or not; if it does, the general integer solution is determined.

## Input

A linear equation $a_{1} x_{1}+\ldots+a_{n} x_{n}=b$, with $a_{i}, b \in \mathbb{Z}, x_{i}$ being integer number unknowns, $i=\overline{1, n}$, and not all $a_{i}=0$.

## Output

Decision on the integer solution of this equation; and if the equation has solutions in $\mathbb{Z}$, its general solution is obtained.

## Method

Step 1. Calculate $d=\left(a_{1}, \ldots, a_{n}\right)$.
Step 2. If $d / b$ then "the equation has integer solution"; go on to Step 3. If $d \times b$ then "the equation does not have integer solution"; stop.

Step 3. Consider $h:=1$. If $|d| \neq 1$, divide the equation by $d$; consider $a_{i}:=a_{i} / d, i=\overline{1, n}, b:=b / d$.

Step 4. Calculate $a=\min _{a_{s} \neq 0}\left|a_{s}\right|$ and determine an $i$ such that $a_{i}=a$.
Step 5. If $a \neq 1$ then go to Step 7.
Step 6. If $a=1$, then:
(A) $x_{i}=-\left(a_{1} x_{1}+\ldots+a_{i-1} x_{i-1}+a_{i+1} x_{i+1}+\ldots+a_{n} x_{n}-b\right) \cdot a_{i}$
(B) Substitute the value of $x_{i}$ in the values of the other determined unknowns.
(C) Substitute integer number parameters for all the variables of the unknown values in the right term: $k_{1}, k_{2}, \ldots, k_{n-2}$, and $k_{n-1}$ respectively.
(D) Write, for your records, the general solution thus determined; stop.

Step7. Write down all $a_{j}, j \neq i$ and under the form:

$$
\begin{aligned}
& a_{j}=a_{i} q_{j}+r_{j} \\
& b=a_{i} q+r \text { where } q_{j}=\left[\frac{a_{j}}{a_{i}}\right], q=\left[\frac{b}{a_{i}}\right] .
\end{aligned}
$$

Step 8. Write $x_{i}=-q_{1} x_{1}-\ldots-q_{i-1} x_{i-1}-q_{i+1} x_{i+1}-\ldots-q_{n} x_{n}+q-t_{h}$. Substitute the value of $x_{i}$ in the values of the other determined unknowns.

Step 9. Consider
and go back to Step 4.
Lemma 1. The previous algorithm is finite.
Proof:
Let's $a_{1} x_{1}+\ldots+a_{n} x_{n}=b$ be the initial linear equation, with not all $a_{i}=0$; check for $\min _{a_{s} \neq 0}\left|a_{s}\right|=a_{1} \neq 1$ (if not, it is renumbered). Following the algorithm, once we pass from this initial equation to a new equation: $a_{1}^{\prime} x_{1}+a_{2}^{\prime} x_{2}+\ldots+a_{n}^{\prime} x_{n}=b^{\prime}$, with $\left|a_{1}^{\prime}\right|<\left|a_{i}\right|$ for $i=\overline{2, n},\left|b^{\prime}\right|<|b|$ and $a_{1}^{\prime}=-a_{1}$.

It follows that $\min _{a_{s} \neq 0}\left|a_{s}\right|<\min _{a_{s} \neq 1}\left|a_{s}\right|$. We continue similarly and after a finite number of steps we obtain, at Step $4, a:=1$ (the actual $a$ is always smaller than the previous $a$, according to the previous note) and in this case the algorithm terminates.

Lemma 2. Let the linear equation be:

$$
\begin{equation*}
a_{1} x_{1}+a_{2} x_{2}+\ldots+a_{n} x_{n}=b, \text { with } \min _{a_{s} \neq 0}\left|a_{s}\right|=a_{1} \text { and the equation } \tag{25}
\end{equation*}
$$

$$
\begin{equation*}
-a_{1} t_{1}+r_{2} x_{2}+\ldots+r_{n} x_{n}=r, \quad \text { with } \quad t_{1}=-x_{1}-q_{2} x_{2}-\ldots-q_{n} x_{n}+q, \quad \text { where } \tag{26}
\end{equation*}
$$ $r_{i}=a_{i}-a_{i} q_{i}, \quad i=\overline{2, n}, \quad r=b-a_{1} q \quad$ while $\quad q_{i}=\left[\frac{a_{i}}{a}\right], \quad r=\left[\frac{b}{a_{1}}\right] . \quad$ Then $\quad x_{1}=x_{1}^{0}$, $x_{2}=x_{2}^{0}, \ldots, x_{n}=x_{n}^{0}$ is a particular solution of equation (25) if and only if $t_{1}=t_{1}^{0}=-x_{1}-q_{2} x_{2}^{0}-\ldots-q_{n} x_{n}^{0}+q, \quad x_{2}, \ldots, x_{n}=x_{n}^{0} \quad$ is a particular solution of equation (26).

## Proof:

$x_{1}=x_{1}^{0}, \quad x_{2}=x_{2}^{0}, \ldots, x_{n}=x_{n}^{0}, \quad$ is a particular solution of equation $(25) \Leftrightarrow$ $a_{1} x_{1}^{0}+a_{2} x_{2}^{0}+\ldots+a_{n} x_{n}^{0}=b \Leftrightarrow a_{1} x_{1}^{0}+\left(r_{2}+a_{1} q_{2}\right) x_{2}^{0}+\ldots+\left(r_{n}+a_{1} q_{n}\right) x_{n}^{0}=a_{1} q+r \Leftrightarrow$ $r_{2} x_{2}^{0}+\ldots+r_{n} x_{n}^{0}-a_{1}\left(-x_{1}^{0}-q_{2} x_{2}^{0}-\ldots-q_{n} x_{n}^{0}+q\right)=r \Leftrightarrow-a_{1} t_{1}^{0}+r_{2} x_{2}^{0}+\ldots+r_{n} x_{n}^{0}=r \Leftrightarrow$ $\Leftrightarrow t_{1}=t_{1}^{0}, x_{2}=x_{2}^{0}, \ldots, x_{n}=x_{n}^{0}$ is a particular solution of equation (26).

Lemma 3. $x_{i}=c_{i 1} k_{1}+\ldots+c_{i n-1} k_{n-1}+d_{i}, i=\overline{1, n}$, is the general solution of equation (25) if and only if

$$
\begin{align*}
& t_{1}=-\left(c_{11}+q_{2} c_{21}+\ldots+q_{n} c_{n 1}\right) k_{1}-\ldots-\left(c_{1 n-1}+q_{2} c_{2 n-1}+\ldots+q_{n} c_{n n-1}\right) k_{n}-  \tag{28}\\
& -\left(d_{1}+q_{2} d_{2}+\ldots+q_{n} d_{n}\right)+q, \\
& x_{j}=c_{1 j 1} k_{1}+\ldots+c_{j n-1} k_{n-1}+d_{j}, \quad j=\overline{2, n}
\end{align*}
$$

is a general solution for equation (26).
Proof:

$$
t_{1}=t_{1}^{0}=-x_{1}^{0}-q_{2} x_{2}^{0}-\ldots-q_{n} x_{n}^{0}+q, x_{2}=x_{2}^{0}, \ldots, x_{n}=x_{n}^{0} \text { is a particular solution of }
$$ the equation (25) $\Leftrightarrow x_{1}=x_{1}^{0}, x_{2}=x_{2}^{0}, \ldots, x_{n}=x_{n}^{0}$ is a particular solution of equation (26) $\Leftrightarrow \exists k_{1}=k_{1}^{0} \in \mathbb{Z}, \ldots, k_{n}=k_{n}^{0} \in \mathbb{Z}$ such that

$$
x_{i}=c_{i 1} k_{1}^{0}+\ldots+c_{i n-1} k_{n-1}^{0}+d_{i}=x_{i}^{0}, i=\overline{1, n} \Leftrightarrow \exists k_{1}=k_{1}^{0} \in \mathbb{Z}, \ldots, k_{n}=k_{n}^{0} \in \mathbb{Z}
$$

such that

$$
x_{i}=c_{i 1} k_{1}^{0}+\ldots+c_{i n-1} k_{n-1}^{0}+d_{i}=x_{i}^{0}, \quad i=\overline{2, n}
$$

and

$$
\begin{aligned}
& t_{1}=-\left(c_{11}+q_{2} c_{21}+\ldots+q_{n} c_{n 1}\right) k_{1}^{0}-\ldots-\left(c_{1 n-1}+q_{2} c_{2 n-1}+\ldots+q_{n} c_{n n-1}\right) k_{n-1}^{0}- \\
& \left(d_{1}+q_{2} d_{2}+\ldots+q_{n} d_{n}\right)+q=-x_{1}^{0}-q_{2} x_{2}^{0}-\ldots-q_{n} x_{n}^{0}+q=t_{1}^{0}
\end{aligned}
$$

Lemma 4. The linear equation

$$
\begin{equation*}
a_{1} x_{1}+a_{2} x_{2}+\ldots+a_{n} x_{n}=b \text { with }\left|a_{1}\right|=1 \text { has the general solution: } \tag{29}
\end{equation*}
$$

$$
\left\{\begin{array}{l}
x_{1}=-\left(a_{2} k_{2}+\ldots+a_{n} k_{n}-b\right) a_{1}  \tag{30}\\
x_{i}=k_{i} \in \mathbb{Z} \\
i=\overline{2, n}
\end{array}\right.
$$

## Proof:

Let's consider $x_{1}=x_{1}^{0}, x_{2}=x_{2}^{0}, \ldots, x_{n}=x_{n}^{0}$, a particular solution of equation (29). $\exists k_{2}=x_{2}^{0}, k_{n}=x_{n}^{0}$, such that $x_{1}=\left(-a_{2} x_{2}^{0}+\ldots+a_{n} x_{n}^{0}-b\right) a_{1}=x_{1}^{0}, x_{2}=x_{2}^{0}, \ldots, x_{n}=x_{n}^{0}$.

Lemma 5. Let's consider the linear equation $a_{1} x_{1}+a_{2} x_{2}+\ldots+a_{n} x_{n}=b$, with $\min _{a_{s} \neq 0}\left|a_{s}\right|=a_{1}$ and $a_{i}=a_{1} q_{i}, i=\overline{2, n}$.

Then, the general solution of the equation is:
$\left\{\begin{array}{l}x_{1}=-\left(q_{2} k_{2}+\ldots+q_{n} k_{n}-q\right) \\ x_{i}=k_{i} \in \mathbb{Z} \\ i=\overline{2, n}\end{array}\right.$

## Proof:

Dividing the equation by $a_{1}$ the conditions of Lemma 4 are met.
Theorem of Correctness. The preceding algorithm calculates correctly the general solution of the linear equation $a_{1} x_{1}+\ldots+a_{n} x_{n}=b$, with not all $a_{i}=0$.

## Proof:

The algorithm is finite according to Lemma 1. The correctness of steps 1, 2, and 3 is obvious. At step 4 there is always $\min _{a_{s} \neq 0}\left|a_{s}\right|$ as not all $a_{i}=0$. The correctness of substep 6 A) results from Lemmas 4 and 5 , respectively. This algorithm represents a method of obtaining the general solution of the initial equation by means of the general solutions of the linear equation obtained after the algorithm was followed several times (according
to Lemmas 2 and 3); from Lemma 3, it follows that to obtain the general solution of the initial linear equation is equivalent to calculate the general solution of an equation at step 6 A), equation whose general solution is given in algorithm (according to Lemmas 4 and 5). The Theorem of correctness has been fully proven.

Note. At step 4 of the algorithm we consider $a:=\min _{a_{s} \neq 0}\left|a_{s}\right|$ such that the number of iterations is as small as possible. The algorithm works if we consider $a:=\left|a_{i}\right| \neq \max _{s=1, n}\left|a_{s}\right|$ but it takes longer. The algorithm can be introduced into a computer program.

## Application

Calculate the integer solution of the equation:

$$
6 x_{1}-12 x_{2}-8 x_{3}+22 x_{4}=14 .
$$

## Solution

The previous algorithm is applied.

1. $(6,-12,-8,22)=2$
2. $2 \mid 14$ therefore the solution of the equation is in $\mathbb{Z}$.
3. $h:=1 ;|2| \neq 1$; dividing the equation by 2 we obtain:
$3 x_{1}=6 x_{2}-4 x_{3}+11 x_{4}=7$.
4. $a:=\min \{|3|,|-6|,|-4|,|11|\}=3, i=1$
5. $a \neq 1$
6. $-6=3 \cdot(-2)+0$
$-4=3 \cdot(-2)+2$
$11=3 \cdot 3+2$
$7=3 \cdot 2+1$
7. $x_{1}=2 x_{2}+2 x_{3}-3 x_{4}+2-t_{1}$
8. 

$$
\begin{array}{ll}
a_{2}:=0 & a_{1}:=-3 \\
a_{3}:=2 & b:=1 \\
a_{4}:=2 & x_{1}:=t_{1} \\
& h:=2
\end{array}
$$

4. We have a new equation:

$$
\begin{aligned}
& -3 t_{1}-0 \cdot x_{2}+2 x_{3}+2 x_{4}=1 \\
& a:=\min \{|-3|,|2|,|2|\} \text { and } \\
& i=3
\end{aligned}
$$

5. $a \neq 1$
6. $-3=2 \cdot(-2)+1$
$0=2 \cdot 0+0$
$2=2 \cdot 1+0$
$1=2 \cdot 0+0$
7. $x_{3}=2 t_{1}+0 \cdot x_{2}-x_{4}+0-t_{2}$. Substituting the value of $x_{3}$ in the value determined for $x_{1}$ we obtain: $x_{1}=2 x_{2}-5 x_{4}+3 t_{1}-2 t_{2}+2$
8. $a_{1}:=1 \quad a_{3}:=-2$
$a_{2}:=0 \quad b:=1$
$a_{4}:=0 \quad x_{3}:=t_{2}$
$h:=3$
9. We have obtained the equation:
$1 \cdot t_{2}+0 \cdot x_{2}-2 \cdot t_{2}+0 \cdot x_{4}=1, a=1$, and $i=1$
10. (A) $t_{1}=-\left(0 \cdot x_{2}-2 t_{2}+0 \cdot x_{4}-1\right) \cdot 1=2 t_{2}+1$
(B) Substituting the value of $t_{1}$ in the values of $x_{1}$ and $x_{3}$ previously determined, we obtain:

$$
\begin{aligned}
& x_{1}=2 x_{2}-5 x_{4}+4 t_{2}+5 \text { and } \\
& x_{3}=-x_{4}+3 t_{2}+2
\end{aligned}
$$

(C) $x_{2}:=k_{1}, x_{4}:=k_{2}, t_{2}:=k_{3}, k_{1}, k_{2}, k_{3} \in \mathbb{Z}$
(D) The general solution of the initial equation is:

$$
\begin{aligned}
& x_{1}=2 k_{1}-5 k_{2}+4 k_{3}+5 \\
& x_{2}=k_{1} \\
& x_{3}=-k_{2}+3 k_{3}+2 \\
& x_{4}=k_{2} \\
& k_{1}, k_{2}, k_{3} \text { are parameters } \in \mathbb{Z}
\end{aligned}
$$

## REFERENCE

[1] Smarandache, Florentin - Whole number solution of equations and systems of equations - part of the diploma thesis, University of Craiova, 1979.

