

A CLASS OF RECURSIVE SETS

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In this article one builds a class of recursive sets, one establishes properties of these sets and one proposes applications. This article widens some results of [1].

1) Definitions, properties.

One calls recursive sets the sets of elements which are built in a recursive manner: let T be a set of elements and f_i for i between 1 and s , of operations n_i , such that $f_i : T^{n_i} \rightarrow T$. Let's build by recurrence the set M included in T and such that:

(Def. 1) 1°) certain elements a_1, \dots, a_n of T , belong to M .

2°) if $(\alpha_{i_1}, \dots, \alpha_{i_{n_i}})$ belong to M , then $f_i(\alpha_{i_1}, \dots, \alpha_{i_{n_i}})$ belong to M for all $i \in \{1, 2, \dots, s\}$.

3°) each element of M is obtained by applying a number finite of times the rules 1° or 2°.

We will prove several proprieties of these sets M , which will result from the manner in which they were defined. The set M is the representative of a class of recursive sets because in the rules 1° and 2°, by particularizing the elements a_1, \dots, a_n respectively f_1, \dots, f_s one obtains different sets.

Remark 1 : To obtain an element of M , it is necessary to apply initially the rule 1.

(Def. 2) The elements of M are called elements M -recursive.

(Def. 3) One calls order of an element a of M the smallest natural $p \geq 1$ which has the propriety that a is obtained by applying p times the rule 1° or 2°.

One notes M_p the set which contains all the elements of order p of M . It is obvious that $M_1 = \{a_1, \dots, a_n\}$.

$$M_2 = \bigcup_{i=1}^s \left\{ \bigcup_{(\alpha_{i_1}, \dots, \alpha_{i_{n_i}}) \in M_1^{n_i}} f_i(\alpha_{i_1}, \dots, \alpha_{i_{n_i}}) \right\} \setminus M_1.$$

One withdraws M_1 because it is possible that $f_j(a_{j_1}, \dots, a_{j_{n_j}}) = a_i$ which belongs to M_1 , and thus does not belong to M_2 .

One proves that for $k \geq 1$ one has:

$$M_{k+1} = \bigcup_{i=1}^s \left\{ \bigcup_{(\alpha_{i_1}, \dots, \alpha_{i_{n_i}}) \in \prod_k^{(i)}} f_i(\alpha_{i_1}, \dots, \alpha_{i_{n_i}}) \right\} \setminus \bigcup_{h=1}^k M_h$$

where each

$\prod_k^{(i)} = \left\{ (\alpha_{i_1}, \dots, \alpha_{i_{n_i}}) / \alpha_{i_j} \in M_{q_j} \quad j \in \{1, 2, \dots, n_i\}; 1 \leq q_j \leq k \text{ and at least an element } a_{i_{j_0}} \in M_{q_j}, 1 \leq j_0 \leq n_i \right\}$.

The sets M_p , $p \in \infty^*$ form a partition of the set M .

Theorem 1:

$$M = \bigcup_{p \in \infty^*} M_p, \text{ where } \infty^* = \{1, 2, 3, \dots\}.$$

Proof:

From the rule 1° it results that $M_1 \subseteq M$.

One supposes that this propriety is true for values which are less than p . It results that $M_p \subseteq M$, because M_p is obtained by applying the rule 2° to the elements of $\bigcup_{i=1}^{p-1} M_i$.

Thus $\bigcup_{p \in \infty^*} M_p \subseteq M$. Reciprocally, one has the inclusion in the contrary sense in accordance with the rule 3°.

Theorem 2: The set M is the smallest set, which has the properties 1° and 2°.

Proof:

Let R be the smallest set having properties 1° and 2°. One will prove that this set is unique.

Let's suppose that there exists another set R' having properties 1° and 2°, which is the smallest. Because R is the smallest set having these proprieties, and because R' has these proprieties also, it results that $R \subseteq R'$; of an analogue manner, we have $R' \subseteq R$: therefore $R = R'$.

It is evident that $M' \subseteq R$. One supposes that $M_i \subseteq R$ for $1 \leq i < p$. Then (rule 3°), and taking in consideration the fact that each element of M_p is obtained by applying rule 2° to certain elements of M_i , $1 \leq i < p$ it results that $M_p \subseteq R$. Therefore $\bigcup_p M_p \subseteq R$ ($p \in \infty^*$), thus $M \subseteq R$. And because R is unique, $M = R$.

Remark 2. The theorem 2 replaces the rule 3° of the recursive definition of the set M by: " M is the smallest set that satisfies proprieties 1° and 2°".

Theorem 3: M is the intersection of all the sets of T which satisfy conditions 1° and 2°.

Proof:

Let T_{12} be the family of all sets of T satisfying the conditions 1^o and 2^o. We note $I = \bigcap_{A \in T_{12}} A$.

I has the properties 1^o and 2^o because:

- 1) For all $i \in \{1, 2, \dots, n\}$, $a_i \in I$, because $a_i \in A$ for all A of T_{12} .
- 2) If $\alpha_{i_1}, \dots, \alpha_{i_{n_i}} \in I$, it results that $\alpha_{i_1}, \dots, \alpha_{i_{n_i}}$ belong to A that is A of T_{12} .

Therefore,

$$\forall i \in \{1, 2, \dots, s\}, f_i(\alpha_{i_1}, \dots, \alpha_{i_{n_i}}) \in A \text{ which is } A \text{ of } T_{12}, \text{ therefore } f_i(\alpha_{i_1}, \dots, \alpha_{i_{n_i}}) \in I$$

for all i from $\{1, 2, \dots, s\}$.

From theorem 2 it results that $M \subseteq I$.

Because M satisfies the conditions 1^o and 2^o, it results that $M \in T_{12}$, from which $I \subseteq M$. Therefore $M = I$

(Def. 4) A set $A \subseteq I$ is called closed for the operation f_{i_0} if and only if for all $\alpha_{i_0 1}, \dots, \alpha_{i_0 n_{i_0}}$ of A , one has $f_{i_0}(\alpha_{i_0 1}, \dots, \alpha_{i_0 n_{i_0}})$ belong to A .

(Def. 5) A set $A \subseteq T$ is called closed M -recursive if and only if:

- 1) $\{a_1, \dots, a_n\} \subseteq A$.
- 2) A is closed in respect to operations f_1, \dots, f_s .

With these definitions, the precedent theorems become:

Theorem 2': the set M is the smallest closed M -recursive set.

Theorem 3': M is the intersection of all closed M -recursive sets.

(Def. 6) The system of elements $\langle \alpha_1, \dots, \alpha_m \rangle$, $m \geq 1$ and $\alpha_i \in T$ for $i \in \{1, 2, \dots, m\}$, constitute a description M -recursive for the element α , if $\alpha_m = \alpha$ and that each α_i ($i \in \{1, 2, \dots, m\}$) satisfies at least one of the proprieties:

$$1) \alpha_i \in \{a_1, \dots, a_n\}.$$

2) α_i is obtained starting with the elements which precede it in the system by applying the functions f_j , $1 \leq j \leq s$ defined by property 2^o of (Def. 1).

(Def. 7) The number m of this system is called the length of the M -recursive description for the element α .

Remark 3: If the element α admits a M -recursive description, then it admits an infinity of such descriptions.

Indeed, if $\langle \alpha_1, \dots, \alpha_m \rangle$ is a M -recursive description of α then

$\left\langle \underbrace{a_1, \dots, a_1}_{h \text{ times}}, \alpha_1, \dots, \alpha_m \right\rangle$ is also a M -recursive description for α , h being able to take all

values from ∞ .

Theorem 4: The set M is identical with the set of all elements of T which admit a M -recursive description.

Proof: let D be the set of all elements, which admit a M -recursive description. We will prove by recurrence that $M_p \subseteq D$ for all p of ∞^* .

For $p=1$ we have: $M_1 = \{a_1, \dots, a_n\}$, and the a_j , $1 \leq j \leq n$ having as M -recursive description: $\langle a_j \rangle$. Thus $M_1 \subseteq D$. Let's suppose that the property is true for the values smaller than p . M_p is obtained by applying the rule 2° to the elements of

$\bigcup_{i=1}^{p-1} M_i$; $\alpha \in M_p$ implies that $\alpha \in f_j(\alpha_{i_1}, \dots, \alpha_{i_n})$ and $\alpha_{i_j} \in M_{h_j}$ for $h_j < p$ and $1 \leq j \leq n_i$.

But a_{i_j} , $1 \leq j \leq n_i$, admits M -recursive descriptions according to the hypothesis of recurrence, let's have $\langle \beta_{j_1}, \dots, \beta_{j_{s_j}} \rangle$. Then $\langle \beta_{1_1}, \dots, \beta_{1_{s_1}}, \beta_{2_1}, \dots, \beta_{2_{s_2}}, \dots, \beta_{n_i 1}, \dots, \beta_{n_i s_{n_i}}, \alpha \rangle$ constitute a M -recursive description for the element α . Therefore if α belongs to D , then $M_p \subseteq D$ which is $M = \bigcup_{p \in \infty^*} M_p \subseteq D$.

Reciprocally, let x belong to D . It admits a M -recursive description $\langle b_1, \dots, b_t \rangle$ with $b_t = x$. It results by recurrence by the length of the M -recursive description of the element x , that $x \in M$. For $t=1$ we have $\langle b_1 \rangle$, $b_1 = x$ and $b_1 \in \{a_1, \dots, a_n\} \subseteq M$. One supposes that all elements y of D which admit a M -recursive description of a length inferior to t belong to M . Let $x \in D$ be described by a system of length t : $\langle b_1, \dots, b_t \rangle$, $b_t = x$. Then $x \in \{a_1, \dots, a_n\} \subseteq M$, where x is obtained by applying the rule 2° to the elements which precede it in the system: b_1, \dots, b_{t-1} . But these elements admit the M -recursive descriptions of length which is smaller than t : $\langle b_1 \rangle, \langle b_1, b_2 \rangle, \dots, \langle b_1, \dots, b_{t-1} \rangle$. According to the hypothesis of the recurrence, b_1, \dots, b_{t-1} belong to M . Therefore b_t belongs also to M . It results that $M \equiv D$.

Theorem 5: Let b_1, \dots, b_q be elements of T , which are obtained from the elements a_1, \dots, a_n by applying a finite number of times the operations f_1, f_2, \dots , or f_s . Then M can be defined recursively in the following mode:

- 1) Certain elements $a_1, \dots, a_n, b_1, \dots, b_q$ of T belong to M .
- 2) M is closed for the applications f_i , with $i \in \{1, 2, \dots, s\}$.
- 3) Each element of M is obtained by applying a finite number of times the rules (1) or (2) which precede.

Proof: evident. Because b_1, \dots, b_q belong to T , and are obtained starting with the elements a_1, \dots, a_n of M by applying a finite number of times the operations f_i , it results that b_1, \dots, b_q belong to M .

Theorem 6: Let's have g_j , $1 \leq j \leq r$, of the operations n_j , where $g_j : T^{n_j} \rightarrow T$ such that M to be closed in rapport to these operations. Then M can be recursively defined in the following manner:

- 1) Certain elements a_1, \dots, a_n de T belong to M .
- 2) M is closed for the operations f_i , $i \in \{1, 2, \dots, s\}$ and g_j , $j \in \{1, 2, \dots, r\}$.
- 3) Each element of M is obtained by applying a finite number of times the precedent rules.

Proof is simple: Because M is closed for the operations g_j (with $j \in \{1, 2, \dots, r\}$), one has, that for any $\alpha_{j_1}, \dots, \alpha_{j_{n_j}}$ from M , $g_j(\alpha_{j_1}, \dots, \alpha_{j_{n_j}}) \in M$ for all $j \in \{1, 2, \dots, r\}$.

From the theorems 5 and 6 it results:

Theorem 7: The set M can be recursively defined in the following manner:

- 1) Certain elements $a_1, \dots, a_n, b_1, \dots, b_q$ of T belong to M .
- 2) M is closed for the operations f_i ($i \in \{1, 2, \dots, s\}$) and for the operations g_j ($j \in \{1, 2, \dots, r\}$) previously defined.
- 3) Each element of M is defined by applying a finite number of times the previous 2 rules.

(Def. 8) The operation f_i conserves the property P iff for any elements $\alpha_{i_1}, \dots, \alpha_{i_{n_i}}$ having the property P , $f_i(\alpha_{i_1}, \dots, \alpha_{i_{n_i}})$ has the property P .

Theorem 8 : If a_1, \dots, a_n have the property P , and if the functions f_1, \dots, f_s preserve this property, then all elements of M have the property P .

Proof:

$M = \bigcup_{p \in \omega^*} M_p$. The elements of M_1 have the property P .

Let's suppose that the elements of M_i for $i < p$ have the property P . Then the elements of M_p also have this property because M_p is obtained by applying the operations f_1, f_2, \dots, f_s to the elements of: $\bigcup_{i=1}^{p-1} M_i$, elements which have the property P .

Therefore, for any p of ω , the elements of M_p have the property P .

Thus all elements of M have it.

Corollary 1 : Let's have the property P : "x can be represented in the form $F(x)$ ".

If a_1, \dots, a_n can be represented in the form $F(a_1), \dots$, respectively $F(a_n)$, and if f_1, \dots, f_s maintains the property P , then all elements α of M can be represented in the form $F(\alpha)$.

Remark. One can find more other equivalent def. of M .

2) APPLICATIONS, EXAMPLES.

In applications, certain general notions like: M - recursive element, M -recursive description, M - recursive closed set will be replaced by the attributes which characterize the set M . For example in the theory of recursive functions, one finds notions like: recursive primitive functions, primitive recursive description, primitively recursive closed sets. In this case " M " has been replaced by the attribute "primitive" which characterizes this class of functions, but it can be replaced by the attributes "general", "partial".

By particularizing the rules 1° and 2° of the def. 1, one obtains several interesting sets:

Example 1: (see [2], pp. 120-122, problem 7.97).

Example 2: The set of terms of a sequence defined by a recurring relation constitutes a recursive set.

Let's consider the sequence: $a_{n+k} = f(a_n, a_{n+1}, \dots, a_{n+k-1})$ for all n of ∞^* , with $a_i = a_i^0$, $1 \leq i \leq k$. One will recursively construct the set $A = \{a_m\}_{m \in \infty^*}$ and one will define in the same time the position of an element in the set A :

1°) a_1^0, \dots, a_k^0 belong to A , and each a_i^0 ($1 \leq i \leq k$) occupies the position i in the set A ;

2°) if $a_n, a_{n+1}, \dots, a_{n+k-1}$ belong to A , and each a_j for $n \leq j \leq n+k-1$ occupies the position j in the set A , then $f(a_n, a_{n+1}, \dots, a_{n+k-1})$ belongs to A and occupies the position $n+k$ in the set A .

3°) each element of B is obtained by applying a finite number of times the rules 1° or 2°.

Example 3: Let $G = \{e, a^1, a^2, \dots, a^p\}$ be a cyclic group generated by the element a . Then G , \square can be recursively defined in the following manner:

1°) a belongs to G .

2°) if b and c belong to G then $b \square c$ belongs to G .

3°) each element of G is obtained by applying a finite number of times the rules 1 or 2.

Example 4: Each finite set $ML = \{x_1, x_2, \dots, x_n\}$ can be recursively defined (with $ML \subseteq T$):

1°) The elements x_1, x_2, \dots, x_n of T belong to ML .

2°) If a belongs to ML , then $f(a)$ belongs to ML , where $f: T \rightarrow T$ such that $f(x) = x$;

3°) Each element of ML is obtained by applying a finite number of times the rules 1° or 2°.

Example 5: Let L be a vectorial space on the commutative corps K and $\{x_1, \dots, x_m\}$ be a base of L . Then L , can be recursively defined in the following manner:

1°) x_1, \dots, x_m belong to L ;

2°) if x, y belong to L and if a belongs to K , then $x \perp y$ y belong to L and $a * x$ belongs to L ;

3°) each element of L is recursively obtained by applying a finite number of times the rules 1° or 2°.

(The operators \perp and $*$ are respectively the internal and external operators of the vectorial space L).

Example 6: Let X be an A -module, and $M \subset X$ ($M \neq \emptyset$), with $M = \{x_i\}_{i \in I}$. The sub-module generated by M is:

$$\langle M \rangle = \{x \in X / x = a_1 x_1 + \dots + a_n x_n, a_i \in A, x_i \in M, i \in 1, \dots, n\}$$

can recursively defined in the following way:

1°) for all i of $\{1, 2, \dots, n\}$, $x_i \in \langle M \rangle$;

2°) if x and y belong to $\langle M \rangle$ and a belongs to A , then $x + y$ belongs to $\langle M \rangle$, and ax also;

3°) each element of $\langle M \rangle$ is obtained by applying a finite number of times the rules 1° or 2°.

In accordance to the paragraph 1 of this article, $\langle M \rangle$ is the smallest sub-set of X that verifies the conditions 1° and 2°, that is $\langle M \rangle$ is the smallest sub-module of X that includes M . $\langle M \rangle$ is also the intersection of all the subsets of X that verify the conditions 1° and 2°, that is $\langle M \rangle$ is the intersection of all sub-modules of X that contain M . One also directly refines some classic results from algebra.

One can also talk about sub-groups or ideal generated by a set: one also obtains some important applications in algebra.

Example 7: One also obtains like an application the theory of formal languages, because, like it was mentioned, each regular language (linear at right) is a regular set and reciprocally. But a regular set on an alphabet $\Sigma = \{a_1, \dots, a_n\}$ can be recursively defined in the following way:

1°) $\emptyset, \{\varepsilon\}, \{a_1\}, \dots, \{a_n\}$ belong to R .

2°) if P and Q belong to R , then $P \cup Q$, PQ , and P^* belong to R , with

$$P \cup Q = \{x / x \in P \text{ or } x \in Q\}; \quad PQ = \{xy / x \in P \text{ and } y \in Q\}, \quad \text{and} \quad P^* = \bigcup_{n=0}^{\infty} P^n \quad \text{with}$$

$$P^n = \underbrace{P \cdot P \cdots P}_{n \text{ times}} \text{ and, by convention, } P^0 = \{\varepsilon\}.$$

3°) Nothing else belongs to R other than those which are obtained by using 1° or 2°.

From which many properties of this class of languages with applications to the programming languages will result.

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