△ –Synchronization Of Interval Nutrosophic Automata

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Abstract

The purpose of this paper is to study Δ –Synchronization of interval neutrosophic automata and their characterizations.

Key words: Interval neutrosophic automaton (INA), Δ –Synchronization.

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1 Introduction

The neutrosophic set was introduced by Florentin Smarandache in 1999 [6]. Fuzzy sets was introduced by Zadeh in 1965 [8]. Bipolar fuzzy sets, YinYang, bipolar fuzzy sets, NPN fuzzy set were introduced by W. R. Zhang in [9, 10, 11]. A NS N is classified by a Truth T_N , Indeterminacy I_N , and Falsity membership F_N where T_N , I_N , and I_N are real standard and non-standard subsets of I_N are subsets of I_N . Fuzzy automaton was introduced by Wee [7]. The INA was introduced by Tahir Mahmood [4]. Retrievability, subsystem, and strong subsystems of INA are studied in the papers I_N , I_N ,

2 Preliminaries

2.1 Definition [5]

A FA is triple F = (T, I, S) where T, I are set of states, set of input symbols and S is transition function in $T \times I \times T \rightarrow [0, 1]$.

2.2 Definition [4]

Let U be universal set. A NS S in U is classified as truth K_s , an indeterminacy L_s and a falsity values M_S where K_s , L_s , and M_S are real standard or non- standard subsets of $]0^- 1^+[$. S = { $\langle z, (K_s(z), L_s(z), M_S(z)), z \in U, K_s, L_s M_S \in]0^- 1^+[$ } and $0^- \le \sup K_s(z) + \sup L_s(z) + \sup M_S(z) \le 3^+$. We take values [0, 1] instead of

 $0^- \le \sup K_s(z) + \sup L_s(z) + \sup M_s(z) \le 3^+$. We take values [0, 1] instead of $]0^-, 1^+[$.

2.3 Definition [4]

Let F = (T, IS) be INA. T and I are set of states and input symbols respectively, and $S = \{(K_s(z), L_s(z), M_s(z))\}$ is an INS in $T \times I \times T$. The set of all strings I is denote by I^* . The empty string is denoted by ϵ and the length of $z \in I^*$ is denoted by |z|.

2.4 Definition [4]

Let F = (T, IS) be INA. Define an INS $s^* = \{(K_{s_*}(z), L_{s_*}(z), M_{S_*}(z))\}$ in $T^* \times I^* \times T$ by

$$K_{s}(\mathbf{z})(t_{a}, \epsilon, t_{b}) = \begin{cases} [1,1] \ if \ t_{a} = t_{b} \\ [0,0] \ if \ t_{a} \neq t_{b} \end{cases}, \quad L_{s}(\mathbf{z})(t_{a}, \epsilon, t_{b}) = \begin{cases} [0,0] \ if \ t_{a} = t_{b} \\ [1,1] \ if \ t_{a} \neq t_{b} \end{cases}, \text{ and }$$

$$M_s(\mathbf{z})(t_a, \epsilon, t_b) = \begin{cases} [0,0] & \text{if } t_a = t_b \\ [1,1] & \text{if } t_a \neq t_b \end{cases}$$

$$K_{s*}(t_a, \mathsf{z}\mathsf{z}', t_b) = \bigvee_{t_r \in T} \left[K_{s*}(t_a, \mathsf{z}, t_r) \land K_{s*}(t_r, \mathsf{z}', t_b) \right] > [0, 0]$$

$$L_{s*}(t_a, zz', t_b) = \Lambda_{t_r \in T} [L_{s*}(t_a, z, t_r) \lor L_{s*}(t_r, z', t_b)] < [1, 1]$$

$$M_{S*}(t_a, zz', t_b) = \Lambda_{t_r \in T} [M_{S*}(t_a, z, t_r) \lor M_{S*}(t_r, z', t_b)] < [1, 1] \forall t_a, t_b \in T, z \in I^* \text{ and } z' \in I.$$

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3 \(\Delta - \Synchronization of Interval Neutrosophic Automata \)

3.1 Definition

Let F = (T, I, S) be an IVNA. F is called deterministic IVNA, $\forall t_a \in T$ and $z \in I$ \exists unique state t_b such that $K_{s*}(t_a, z, t_b) > [0, 0]$, $L_{s*}(t_a, z, t_b) < [1, 1]$, $M_{s*}(t_a, z, t_b) < [1, 1]$.

3.2 Definition

Let F = (T, I, S) be an IVNA and $\Theta = T_1, T_2, \dots, T_z$ be a partition of T. If $K_{s*}(t_a, z, t_b) > [0, 0]$, $L_{s*}(t_a, z, t_b) < [1, 1]$, $M_{s*}(t_a, z, t_b) < [1, 1]$ for some $z \in I$ then $t_a \in T_S$ and $t_b \in T_{S+1}$. Then Θ is periodic partition of order $z \ge 2$. An INA F is periodic of period $z \ge 2$ iff $z = Maxcard(\Theta)$, maximum is consider all periodic partitions Θ of F. F has no periodic partition, then F is called aperiodic.

Note.

Throughout this paper we consider aperiodic INA.

3.3 Definition

Let F = (T, I, S) be an IVNA. Two states t_a , t_b interval neutrosophic stability related (INSR) denoted by $t_a \Omega t_b$, for any string $z \in I^*$, $t_k \in T$ such that

$$\begin{split} &K_{S*}(t_a,zz',t_k) > [0,0] \Leftrightarrow K_{S*}(t_b,zz',t_k) > [0,0] \\ &L_{S*}(t_a,zz',t_k) < [1,1] \Leftrightarrow L_{S*}(t_b,zz',t_k) < [1,1] \\ &M_{S*}(t_a,zz',t_k) > [0,0] \Leftrightarrow M_{S*}(t_b,zz',t_k) < [1,1] \end{split}$$

3.4 Example

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Let F = (T, I, S) be an IVNA, where \{T = T_1, T_2, T_3, T_4\} I = \{z, z'\} and S are defined as below. (K_s, L_s, M_s)(t_1, z, t_4) = \{[0.3, 0.4], [0.4, 0.5], [0.6, 0.8]\} (K_s, L_s, M_s)(t_1, z', t_2) = \{[0.1, 0.2], [0.3, 0.4], [0.7, 0.8]\} (K_s, L_s, M_s)(t_2, z, t_3) = \{[0.2, 0.3], [0.5, 0.6], [0.8, 0.9]\} (K_s, L_s, M_s)(t_2, z', t_4) = \{[0.7, 0.8], [0.3, 0.4], [0.2, 0.3]\}
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 $(K_s, L_s, M_s)(t_2, z', t_4) = \{[0.7,0.8], [0.3,0.4], [0.2,0.3]\}\$ $(K_s, L_s, M_s)(t_3, z, t_2) = \{[0.6,0.7], [0.4,0.5], [0.3,0.4]\}$

 $(K_s, L_s, M_s)(t_3, z', t_4) = \{[0.5, 0.6], [0.4, 0.5], [0.2, 0.3]\}\$ $(K_s, L_s, M_s)(t_4, z, t_1) = \{[0.8, 0.9], [0.2, 0.3], [0.1, 0.2]\}$

 $(K_s, L_s, M_s)(t_4, z', t_3) = \{[0.3, 0.4], [0.4, 0.5], [0.6, 0.8]\}$

For any string $v \in I^*$, there exists a string $zz'z' \in I^*$ such that

 $K_{s*}(t_1, vzz'z', t_k) > [0, 0] \Leftrightarrow K_{s*}(t_4, vzz'z', t_k) > [0, 0]$

 $L_{s*}(t_1, vzz'z', t_k) < [1, 1] \Leftrightarrow L_{s*}(t_4, vzz'z', t_k) < [1, 1]$

 $M_{s*}(t_1, vzz'z', t_k) < [1, 1] \Leftrightarrow M_{s*}(t_4, vzz'z', t_k) < [1, 1]$ and

 $K_{S*}(t_2, vzz'z', t_l) > [0, 0] \Leftrightarrow K_{S*}(t_3, vzz'z', t_l) > [0, 0]$

 $L_{s*}(t_2, vzz'z', t_l) < [1, 1] \Leftrightarrow L_{s*}(t_3, vzz'z', t_l) < [1, 1]$

 $M_{s*}(t_2, vzz'z', t_l) < [1, 1] \Leftrightarrow K_{s*}(t_3, vzz'z', t_l) < [1, 1].$

The states t_1 , t_4 and t_2 , t_3 are interval neutrosophic stability related.

3.5 Definition

Let F = (T, I, S) be an IVNA. F is called Δ –Synchronization if \exists a string $z \in I^*$, $t_b \in T$ and a real number Δ with $\Delta \in (0,1]$ such that $K_{s*}(t_a, z, t_b) \geq \Delta > [0,0]$, $L_{s*}(t_a, z, t_b) \leq \Delta < [1,1]$, $M_{s*}(t_a, z, t_b) \leq \Delta < [1,1]$ \forall $t_a \in T$.

4 Algorithm

Let F = (T, I, S) be an IVNA.

- 1) Find the equivalence classes of the states T using INSR.
- 2) Construct the quotient INA G by considering each equivalence class as a state.
- 3) Relabel the quotient INA along with neutrosophic values G into G' keeping the stability class.
- 4) Construct New INA F' from G'.
- 5) INA G' gives the synchronized string.

4.1 Example

From Example 3.4 and the quotient *INA G* is as follows.

$$\begin{aligned} &(K_{s*},L_{s*},M_{s*})(t_1t_4,z,t_1t_4) = \{[0.3,0.4],[0.4,0.5],[0.6,0.8]\} \\ &(K_{s*},L_{s*},M_{s*})(t_1t_4,z',t_2t_3) = \{[0.1,0.2],[0.4,0.5],[0.7,0.8]\} \\ &(K_{s*},L_{s*},M_{s*})(t_2t_3,z,t_2t_3) = \{[0.2,0.3],[0.5,0.6],[0.8,0.9]\} \\ &(K_{s*},L_{s*},M_{s*})(t_2t_3,z',t_1t_4) = \{[0.5,0.6],[0.4,0.5],[0.2,0.3]\} \end{aligned}$$

Relabled quotient INA G' is as follows

$$\begin{aligned} &(K_{s*},L_{s*},M_{s*})(t_1t_4,z',t_1t_4) = \{[0.1,0.2],[0.4,0.5],[0.7,0.8]\} \\ &(K_{s*},L_{s*},M_{s*})(t_1t_4,z,t_2t_3) = \{[0.3,0.4],[0.4,0.5],[0.6,0.8]\} \\ &(K_{s*},L_{s*},M_{s*})(t_2t_3,z,t_2t_3) = \{[0.2,0.3],[0.5,0.6],[0.8,0.9]\} \\ &(K_{s*},L_{s*},M_{s*})(t_2t_3,z',t_1t_4) = \{[0.5,0.6],[0.4,0.5],[0.2,0.3]\} \end{aligned}$$

Relabled INA F' from G' is as follows

$$\begin{aligned} &(K_s,L_s,M_s)(t_1,z',t_4) = \{[0.5,0.6],[0.4,0.5],[0.2,0.3]\} \\ &(K_s,L_s,M_s)(t_1,z,t_2) = \{[0.3,0.4],[0.4,0.5],[0.6,0.8]\} \\ &(K_s,L_s,M_s)(t_2,z,t_3) = \{[0.2,0.3],[0.5,0.6],[0.8,0.9]\} \\ &(K_s,L_s,M_s)(t_2,z',t_4) = \{[0.7,0.8],[0.3,0.4],[0.2,0.3]\} \\ &(K_s,L_s,M_s)(t_3,z,t_2) = \{[0.6,0.7],[0.4,0.5],[0.3,0.4]\} \\ &(K_s,L_s,M_s)(t_3,z',t_4) = \{[0.8,0.9],[0.2,0.3],[0.1,0.2]\} \\ &(K_s,L_s,M_s)(t_4,z,t_3) = \{[0.3,0.4],[0.4,0.5],[0.6,0.8]\} \\ &(K_s,L_s,M_s)(t_4,z',t_1) = \{[0.3,0.4],[0.4,0.5],[0.6,0.8]\} \end{aligned}$$

In the relabeled INA there exists a string $zz' \in I^*$ in F' such that

$$K_{S*}(t_i, zz', t_4) > [0, 0], L_{S*}(t_i, zz', t_4) < [1, 1] \text{ and } M_{S*}(t_i, zz', t_4) < [1, 1] \ \forall \ t_i \in T.$$

5. Procedure for finding Δ –Synchronized String of Interval Neutrosophic Automata

Let F = (T, I, S) be an INA. We define another INA as follows:

$$F_S = (2^T, I, M_S, T, D \subseteq T)$$
 where

T- Starting state on F_S , D- set of all final states on F_S , M_S — Interval neutrosophic transition function and is defined by

$$K_{M_S}(T_a, z, T_b) = \Lambda \{ (K_S(t_a, z, t_b)) \} > [0, 0]$$

$$L_{M_s}(T_a, z, T_b) = \forall \{(L_s(t_a, z, t_b))\} < [1, 1]$$

$$M_{M_S}(T_a, z, T_b) = \bigvee \{ (M_S(t_a, z, t_b)) < [1, 1], t_a \in T_a, t_b \in T_b, T_a, T_b \in 2^T \text{ for } z \in I.$$

 M_S is a deterministic INA and a string $z \in I$ is Δ – synchronized in F iff \exists a singleton subsets $T_t \in 2^T$ such that

$$K_{M_{S_*}}(T_a, z, T_t) > [0, 0], L_{M_{S_*}}(T_a, z, T_t) < [1, 1] \text{ and } M_{M_{S_*}}(T_a, z, T_b) < [1, 1].$$

6 Conclusion

 Δ –Synchronization of INA are introduce, algorithm is given for finding Synchronized string using interval neutrosophic stability relation. Finally procedure is given for finding synchronized string.

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